# CPSC 416 Distributed Systems

### Winter 2022 Term 2 (March 23, 2023)

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# Logistics



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### **Deadlines**

Project 4 Released. Late Due: April 13, 2023.

Project 5 Released Due: April 13, 2023. No extensions.



All project work is due April 13, 2023. Late projects are scaled to 75% of the on-time max.

Final Exam: April 20, 2023, DMP 310, 08:30-11:00. Format TBA.

### Deadlines

#### Alternate Path 1 & 2: Review in progress

- Piazza private threads need TLC
  - Weekly updates due each Monday @ 23:59 PT
- Final reports due no later than Thursday April 13, 2023 @ 23:59 PT
- Optional 10 min presentation April 13, 2023, up to 10 minutes.

Instructor Office Hours:

- Zoom Office Hours (Tuesday) @ 13:00-14:00
- Discord (Casual) Office Hours (Thursday) @ 14:00-15:00

TA Office Hours:

- Eric: Friday 9-11 am (in-person and Zoom)
- Japraj: Wednesday 3-5 pm (Zoom)
- Yennis: Thursday 2-4 (Zoom), Friday 2-4 (in-person)



# Readings

Required:

Recommended:







### **Questions?**

Questions about the class?

Questions about the previous lecture?

Funny stories to share?



# **Today's Failure**



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## **Github.com Outage**

Event: October 21, 2018 22:52 UTC

Planned outage: goal is to replace a failing 100Gb/s optical network device.

"Connectivity between these two locations was restored in 43 seconds, but this brief outage triggered a chain of events that led to 24 hours and 11 minutes of service degradation."

Infrastruture: MySQL with Orchestrator to manage cluster topologies.

Note: Orchestrator uses **Raft** for consensus.



### **Github.com Outage**

Network goes out: Raft starts "leadership deselection"

Note: optical link was between two Eastern US sites.

West coast data center and East coast Orchestrator form quorum

Fail over to clusters in West coast data center: write operations begin working.

Network fixed: traffic starts going to West coast site

Note: East coast had some updates that had not propagated to west coast yet. This **blocked** primary returning to East coast.



### **Github.com**

Things come unraveled due to increased latency, unexpected topologies. Decision to degrade service rather than compromise consistency.

Start restoring databases from backup.

Restoration was started October 22, 2018 00:05 UTC Restoration completed and service restored October 22, 2018 23:03 UTC

Twenty three hours to restore from a 43 second network disruption.

Takeaway: Recovery is the hard part.

<u>Source</u>



# **Lesson Goals**



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### **Distributed Data Analytics**

MapReduce

Spark (and RDDs)



Data Parallel (Divide & Conquer):

- Divide Data across nodes
- Load balancing, decomposition
- Messaging for data dependencies
- Application usage







#### Pipelining

- Divide work into smaller tasks
  - Small number of tasks per node
  - Faster than generality
- Data streamed in chunks through task pipeline
- Increases throughput



#### Model Parallelism

- Divide state across nodes
- Less processing per node
- Input passed to all nodes
- Output combined from all nodes
- Must handle dependencies





#### distribute input



#### Model Parallelism

- Divide state across nodes
- Less processing per node
- UBC
- Input passed to all nodes
- Output combined from all nodes
- Must handle dependencies

## MapReduce

MapReduce: Simplified Data Processing on Large Clusters, J. Dean, OSDI 2004.

- Hadoop MapReduce
- AWS infrastructure



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## MapReduce

#### Input:

• Set of key-value pair records

#### Map function

- Input: unique key-value pair
- Output: a new key-value pair

#### Reduce function:

- Input: output from map function
- Output: final result

Master: orchestrates workers, I/O, failure management





## MapReduce



Wordcount example:

• Input: Collection of files

Map function:

- Input: File, key=filename, content=value
- Output: file with key=word, value=list of counts

Reduce function:

- Input: file with key=word, value=list of counts
- Output: list of words with total counts

Other examples:

• URL access frequency, page rank, inverted word index



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### MapReduce

Combining Techniques:

- Data parallel: chunks to mappers
- Pipelining: mapper to reducer
- Model parallelism: reducers process parts of key space, combine

Dataflow model means flow of data determines execution





# **Map Reduce: Design Decisions**

Master data structures:

Tracking

Locality:

• Scheduling, placement of intermediate data

Task granularity:

- Finer granularity: more flexibility, management operation execution time
- Coarse granularity: lower management overhead

Fault tolerance:

- Master: standby replication
- Worker: detect failures or stragglers and re-execute

Failure semantics:

Importance of Consistency and complete results

Backup tasks:

Inevitable failures: speculative task backup





### **Map Reduce: Limitations**



UBC

Failure inevitable: cannot re-execute entire operation

Fault-tolerant mechanism: requires intermediate data availability

- Serialiation to/from persistent storage
- Remote access and data movement

Data amplification:

- Intermediate data may be much larger than input
- Executions are iterative
- Storage level replication

System scale: cannot assume best-in-class storage devices



Resilient Distributed Datasets: A Fault-Tolerant Abstraction for In-Memory Cluster Computing

Faster analytics (10x) versus Hadoop

- Workloads: graph, streaming, SQL, Machine Learning, etc.
- Languages: Java, Python, Scala, etc.
- Platforms: AWS, Kubernetes, etc.

Apache Spark



## **Spark: Goals**

#### Allow in-memory data sharing

- Fast DRAM versus slow hard disk
- No serialization cost

Fault-tolerant







### **Resilient Distributed Datasets: Introduction**



Input Just recompute from storage or other RDDs in lineage Immutable partitioned record collection

#### Created using transformations

- Operations on data in stable storage
- Map/join/filter on other RDDs

Used via actions (count, collect, save) RDDs map back to source

Compute partitions from data in stable storage

Users control persistence and partitioning



### **Resilient Distributed Datasets: Example**

#### Console log mining example





## **Resilient Distributed Datasets: Transformations & Actions**

Transformations	$\begin{array}{c} map(f: T \Rightarrow U) \\ filter(f: T \Rightarrow Bool) \\ flatMap(f: T \Rightarrow Seq[U]) \\ sample(fraction : Float) \\ groupByKey() \\ reduceByKey(f: (V, V) \Rightarrow V) \\ union() \\ join() \\ cogroup() \\ crossProduct() \\ mapValues(f: V \Rightarrow W) \\ sort(c: Comparator[K]) \\ p \in Diring(K) \end{array}$	: $RDD[T] \Rightarrow RDD[U]$ : $RDD[T] \Rightarrow RDD[T]$ : $RDD[T] \Rightarrow RDD[T]$ : $RDD[T] \Rightarrow RDD[T]$ (Deterministic sampling) : $RDD[(K, V]] \Rightarrow RDD[(K, Scq[V])]$ : $RDD[(K, V]] \Rightarrow RDD[(K, V)]$ : $(RDD[T], RDD[T]) \Rightarrow RDD[T]$ : $(RDD[(K, V)], RDD[(K, W])] \Rightarrow RDD[(K, (V, W))]$ : $(RDD[(K, V)], RDD[(K, W])] \Rightarrow RDD[(K, (Seq[V], Seq[W]))]$ : $(RDD[T], RDD[U]) \Rightarrow RDD[(T, U]]$ : $(RDD[(K, V)] \Rightarrow RDD[(K, W)]$ (Preserves partitioning) : $RDD[(K, V)] \Rightarrow RDD[(K, V)]$	Narrow Dependencies:	Wide Dependencies:
Actions Table 2: Transf	$count() : count() : count() : collect() : reduce(f : (T, T) \Rightarrow T) : lookup(k : K) : save(path : String) : formations and actions available on$	$\begin{array}{l} \text{RDD}[T] \Rightarrow \text{Long}\\ \text{RDD}[T] \Rightarrow \text{Seq}[T]\\ \text{RDD}[T] \Rightarrow \text{Seq}[T]\\ \text{RDD}[T] \Rightarrow T\\ \text{RDD}[(K, V)] \Rightarrow \text{Seq}[V] (On hash/range partitioned RDDs)\\ \text{Outputs RDD to a storage system, e.g., HDFS}\\ \text{RDDs in Spark. Seq}[T] denotes a sequence of elements of type T. \end{array}$	join with inputs co-partitioned	join with inputs no

### **Resilient Distributed Datasets: Scheduling Action Execution**



Program defines dependencies

Actions:



- Directed acyclic graph (DAG)
- Minimize dependencies
- Optimize parallelism
- Limit I/O contention

Tasks assigned based on data locality

# **Spark: Goals**

Data in memory?

- Distributed shared memory like ٠ runtime
  - Log updates •
  - Persist lineage •



### less data to persist in execution critical path



read data as low as once, less slow storage I/O

Log coarse grained operations applied to all items in RDD elements

more control on locality



## **Spark: Goals**



#### Data in memory?

- Distributed shared memory like runtime
  - Log updates
  - Persist lineage



Log coarse grained operations applied to all items in RDD elements

# **Spark: Evaluation**

#### Up to 20x better than Hadoop

- Iterative
- Machine learning
- Graph applications

Analytics report generation 40x

Rapid failure recovery

1TB dataset queries with 5-7 second latencies





# **Lesson Review**



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## **Distributed Data Analytics**

Systems for scalable data processing

MapReduce

Spark



# **Questions?**



